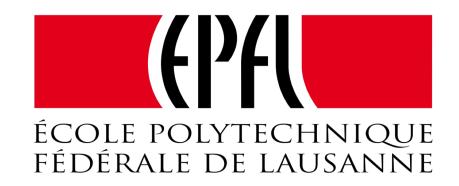
Performance Contracts for Software Network Functions

<u>Rishabh Iyer</u>, Luis Pedrosa, Arseniy Zaostrovnykh, Solal Pirelli, Katerina Argyraki, George Candea



Software Network Functions – Pros and Cons

 \circ Increased flexibility \checkmark

 $_{\odot}$ Reduced capital and operating expenses \checkmark

Programming errors X

Unexpected performance behaviour X

Dealing with unexpected NF performance

Goal: Comprehensive understanding of NF's performance profile

- Operators capacity planning and anticipate attacks
- Developers informed development decisions
- Previous work [NSDI'12, NSDI'18, SIGCOMM'18]
 - Focus on narrow subset of input workloads
 - Offer few completeness guarantees

Performance Contracts for NFs

Abstraction for users to parameterize arbitrary input workloads

• Predict performance for workload spec without running NF

• Performance predicted as function of **Performance Critical Variables (PCVs)**

Per-packet metrics: Instruction count, memory accesses, latency (cycles)

Performance Contract



Outline

 \odot What is a performance contract?

○ How does Bolt generate contracts?

○ Evaluation & Use-Case

void MAC_bridge(pkt* p,port in_port) {

```
if (invalid_hdr(p)) {
    DROP(p);
    return;
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```
}
```

```
MACtable_put(p->src_mac,&in_port);
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```
if (MACtable_get(p->dst_mac,&out_port))
     FORWARD(p,out_port);
also
```

else

```
BROADCAST(p,in_port);
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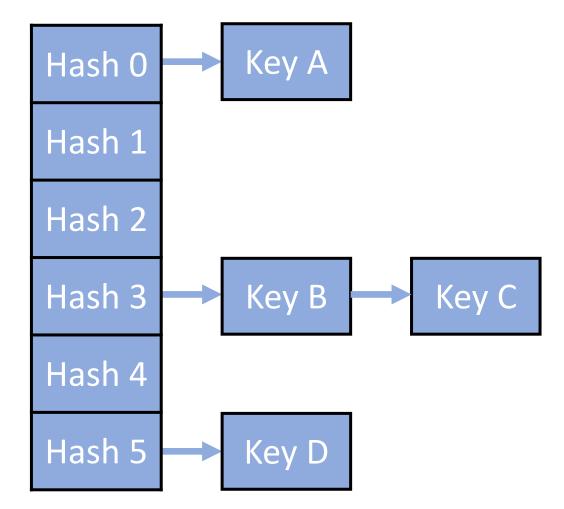
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MACtable implementation



Performance Contracts Example

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Performance Contract for MAC_bridge Metric: Lines of pseudo-code

Traffic Class	Performance

Performance Contracts Example

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Performance Contract for MAC_bridge Metric: Lines of pseudo-code

Traffic Class	Performance
Invalid Header	3
Valid, DestMAC known	3 C + 20
Valid, DestMAC unknown	3 C + 100

Spec 1: Unconstrained traffic

Performance Contract for MAC_bridge Metric: Lines of pseudo-code

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Spec 1: Unconstrained traffic

 $\Rightarrow C = \max_collisions$

Predicted performance: 3(max_collisions) + 100 Performance Contract for MAC_bridge Metric: Lines of pseudo-code

Traffic Class	Performance
Invalid Header	3
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Spec 2: No hash collisions

 $\Rightarrow C = 0$

Predicted performance: 100

Performance Contract for MAC_bridge Metric: Lines of pseudo-code

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DestMAC known

 $\Rightarrow C = 0$

Predicted performance: 20

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Performance Contract for MAC_bridge Metric: Lines of pseudo-code

Traffic Class	Performance
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C = Number of hash collisions

Contracts quantify performance for all traffic classes of the NF Users query contract for performance of specific input workloads

Outline

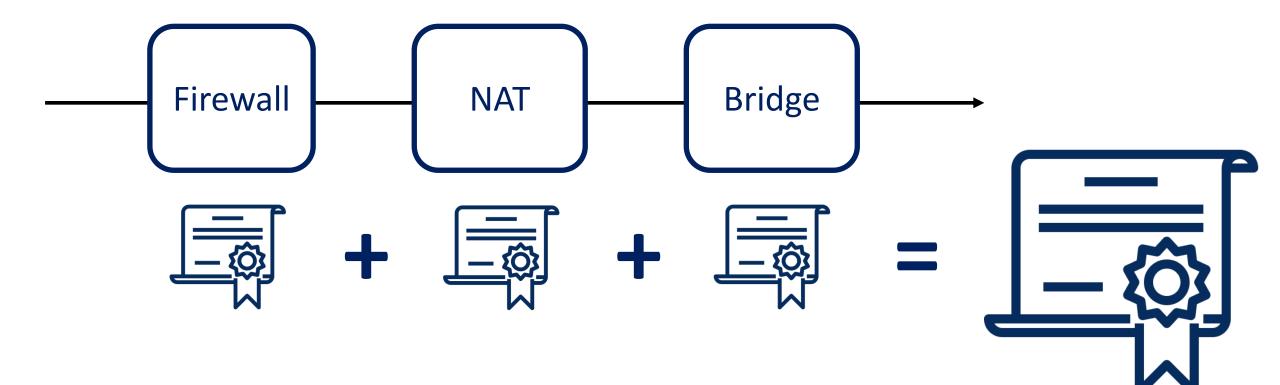
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○ How does Bolt generate contracts?

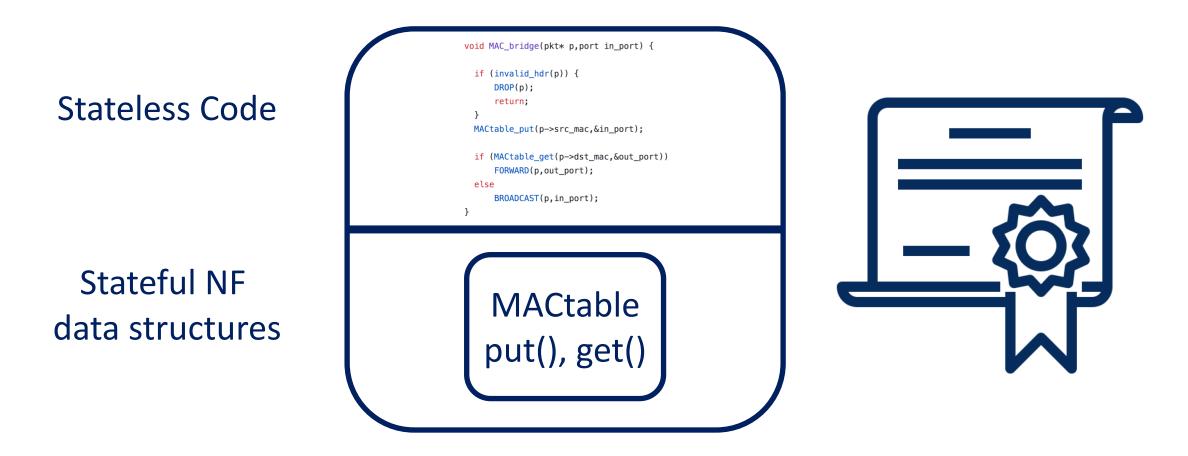
○ Evaluation & Use-Case

Generating performance contracts recursively

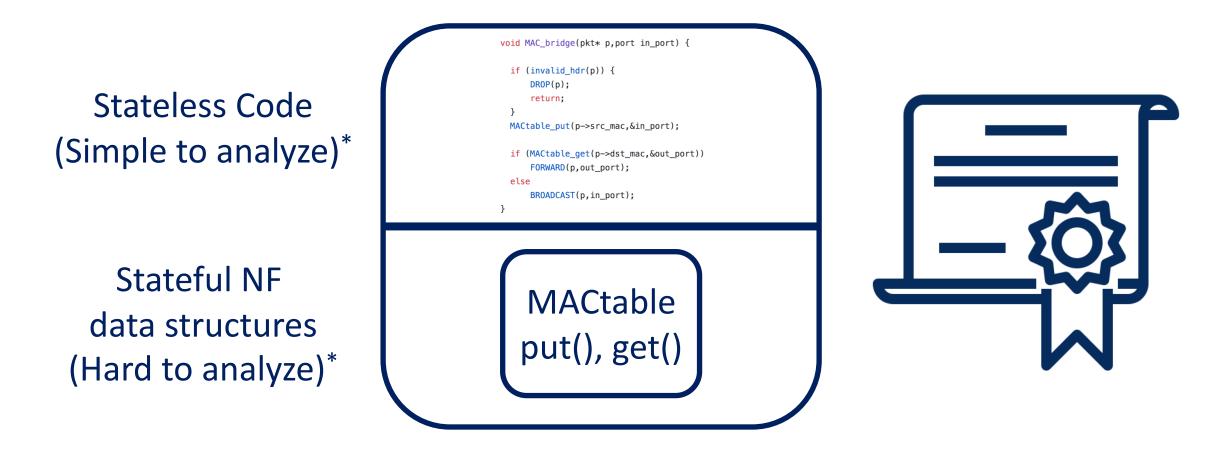
NF chain



Generating performance contracts recursively Individual NF



Generating performance contracts recursively Individual NF



*A.Zaostrovnykh, S.Pirelli, L.Pedrosa, K.Argyraki, G.Candea ''A Formally Verified NAT'' SIGCOMM 2017

Generating performance contracts recursively

 \odot Well defined separation between stateful and stateless NF code*

• NFs typically have well defined, isolated state

Encapsulate NF state using a library of data structures

• Stateful data structures – Base case of recursive process

○ Analyze once, reuse across NFs

*A.Zaostrovnykh, S.Pirelli, L.Pedrosa, K.Argyraki, G.Candea ''A Formally Verified NAT'' SIGCOMM 2017

Analyzing stateful data structures

 $Performance_{NF} = f(input packet, NF state, config, ...)$

Cannot account for all possible packet histories -> Path explosion

• BUT, performance of MACtable depends **ONLY** on number of hash collisions

Performance Critical Variables (PCVs)

○ Abstract away NF state specificities

Succinctly summarize impact of packet history, configuration on performance

 \odot Tailor legibility and detail to audience

Contract for MACtable_put

Traffic Class	Performance
Unconstrained	1 C + 2

Contract for MACtable_get

Traffic Class	Performance
Key present	2 C + 12
Key absent	2 C + 7

C = Number of hash collisions

Only PCV required to summarize perf in terms of lines of pseudo-code

• Symbolically execute stateless code to traverse all execution paths

- While traversing each path
 - Keep track of performance metrics for stateless code
 - Plug in contracts for stateful code using path constraints

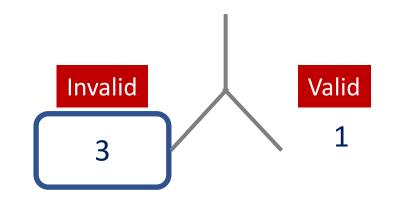
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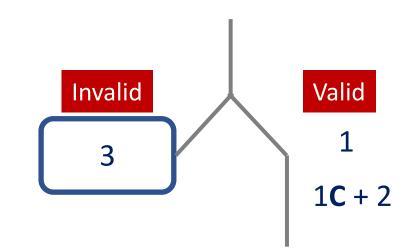
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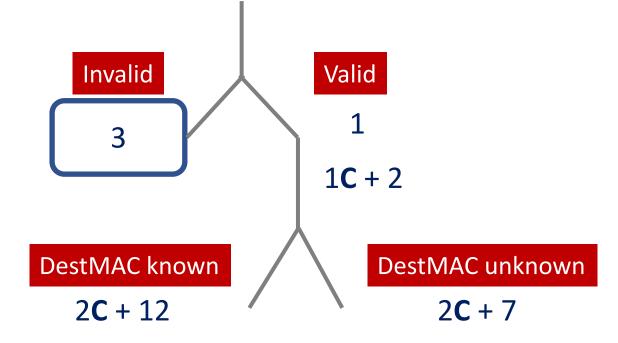
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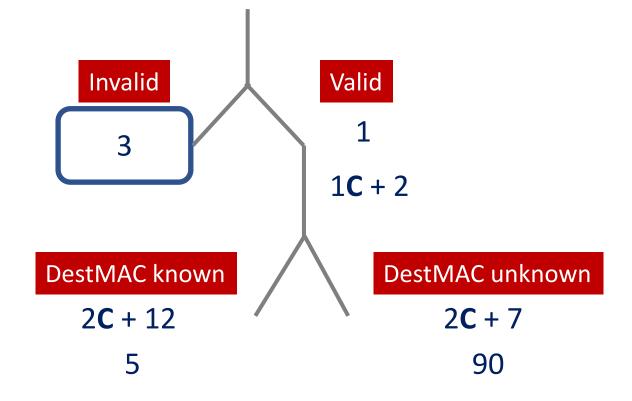
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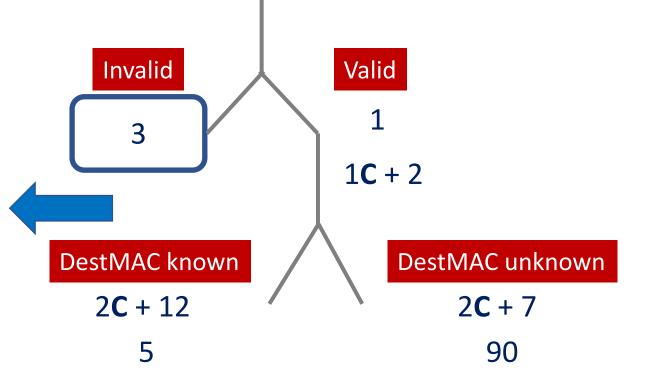
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Performance Contract for MAC_bridge

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Performance Contracts for NF chains

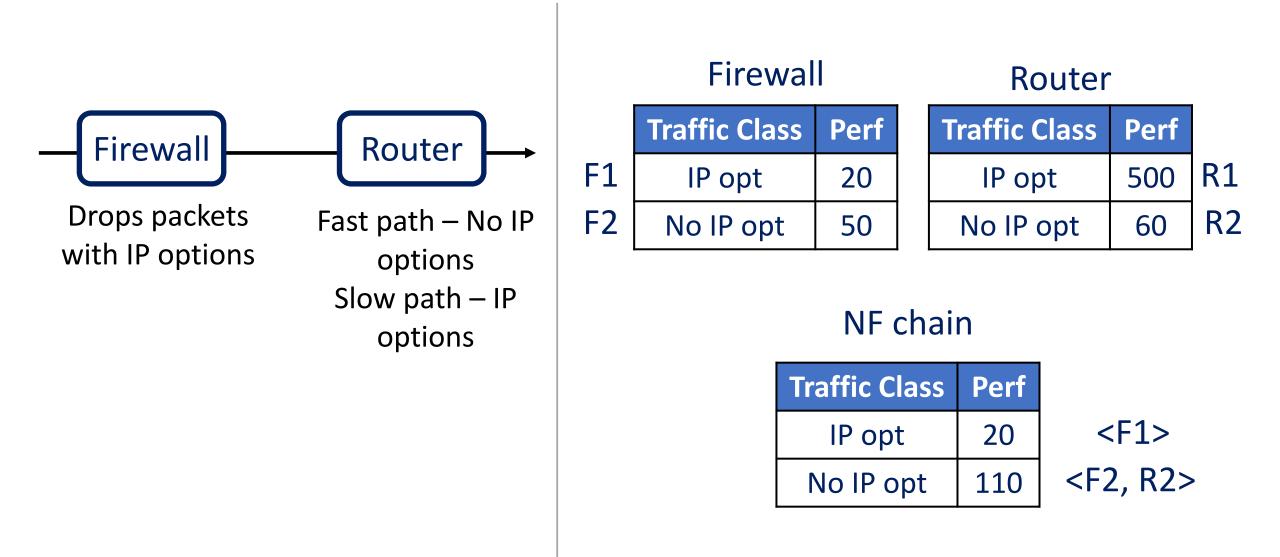
Generate performance contracts for individual NFs in chain

○ Pair together traffic classes from communicating NFs

○ For each pair - AND respective constraints together

Equate packet sent by first NF to packet received by second

Performance Contract for NF chains - Example



Outline

 \odot What is a performance contract?

• How does Bolt generate contracts?

○ Evaluation and Use-Case

Evaluation setup & methodology

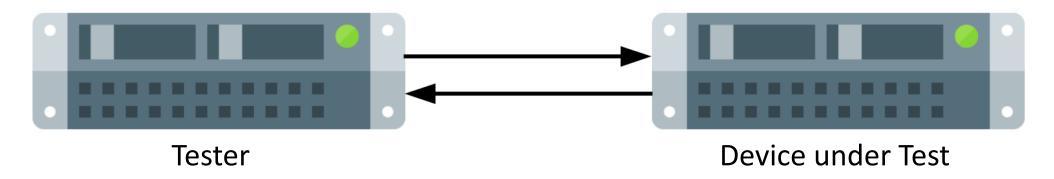
○ 4 NFs - NAT, Maglev-like LB, MAC bridge, LPM router

Analyze NF logic + DPDK + NIC driver*

Metrics – <u>instructions executed</u>, <u>memory accesses</u>, execution cycles

Testbed - Intel Xeon E5-2667v2 3.3GHz, 82599ES 10Gb NICs

Compare predicted vs measured performance for various packet classes

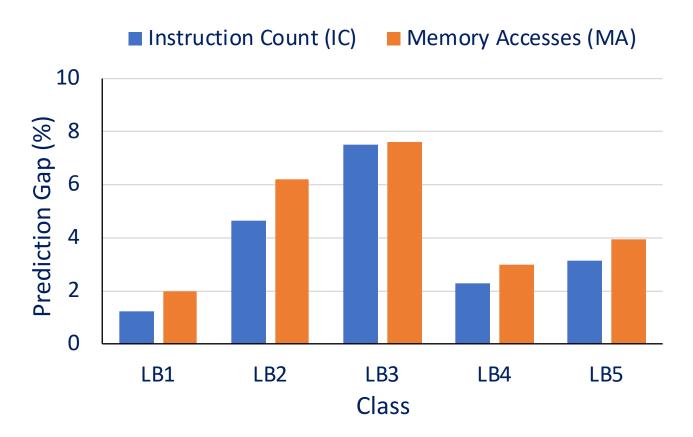


*S.Pirelli, A.Zaostrovnykh, G.Candea ''A Formally Verified NAT Stack'' KBNETS Workshop - SIGCOMM 2018

Predictions for Instruction Count, Memory Accesses

Results for Maglev-like Load Balancer

Class	Description
LB1	Unconstrained traffic
LB2	Client packet, new flow
LB3	Client packet, existing flow, unresponsive backend
LB4	Client packet, existing flow, existing backend
LB5	Heartbeat packets



Max prediction gap – 7.5% (IC) and 7.6% (MA)

Why is there a prediction gap?

Source 1: Trade-off between precision and legibility in PCVs

Can be overcome by exposing more detail

Source 2: Differences between analyzed and production code

Disabled link time optimizations in analyzed code

Use Case – Informed cost-benefit analysis

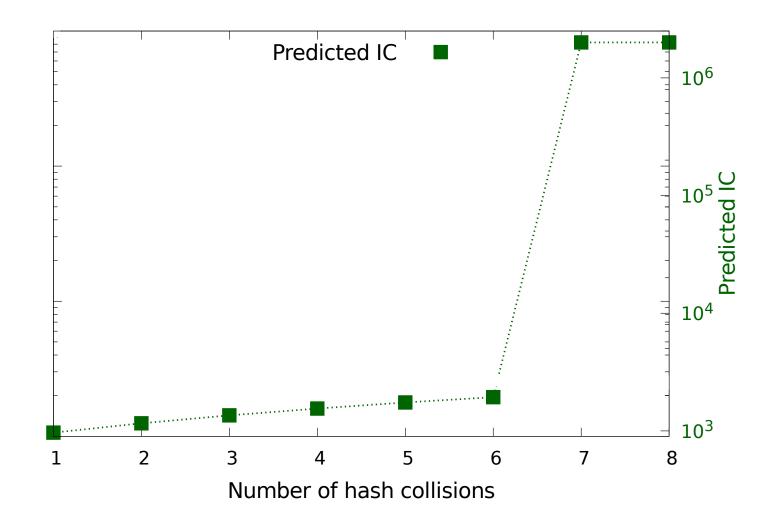
 \odot Example: Bridge with randomized hash table

- Incorporates random key into hash function
- Rehashes all entries with a new key when collisions greater than a threshold

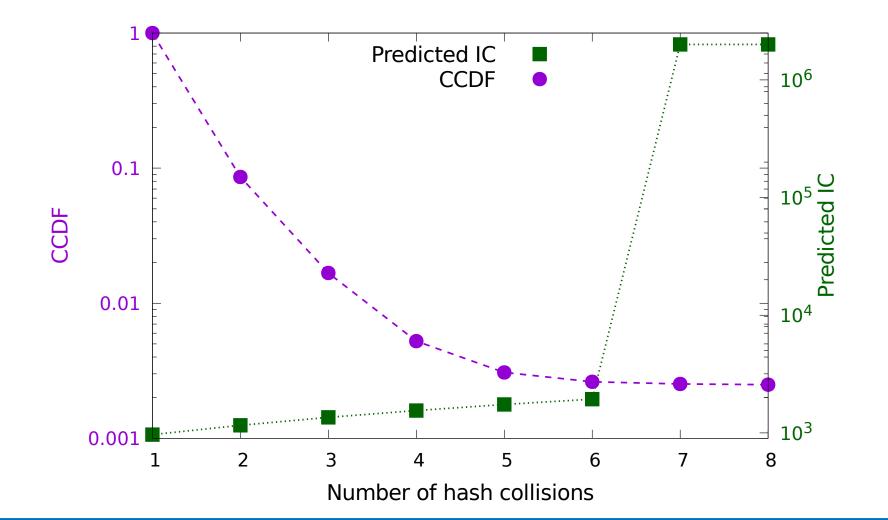
 \odot Question: Where to place threshold?

- Avoid rehashing under normal operation
- Should rehash under attack

Use Case – Informed cost-benefit analysis



Use Case – Informed cost-benefit analysis



Bolt allows operators to visualize the consequences of their decisions

Performance Contracts for NFs

Abstraction for users to parameterize arbitrary input workloads

• Predict performance for workload spec without running NF

• Performance predicted as function of **Performance Critical Variables (PCVs)**



bolt-perf-contracts.github.io

Backup Slides

○ <u>Distiller</u>

- <u>Results IC, MA</u>
- <u>Results NF chains</u>
- <u>Results Latency</u>
- Full Blown Contract

The Bolt Distiller

 \odot Users need to know which traffic classes are likely

• Bolt is a static analysis tool, cannot know probabilities of each traffic class

• The Bolt Distiller

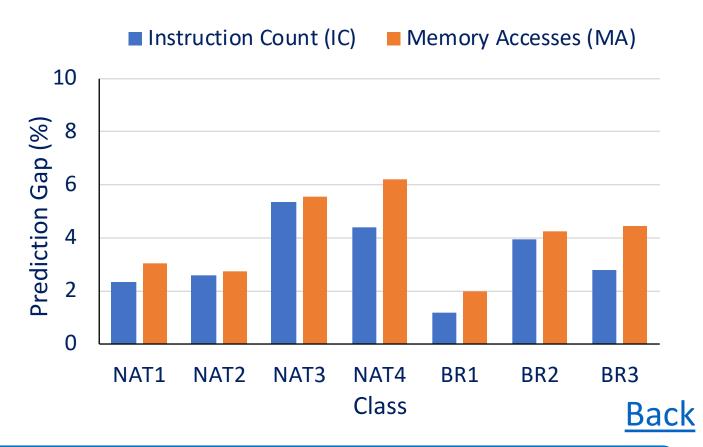
- Input A representative packet trace
- Output Execution path taken by each packet & values of PCVs
- Users can then extrapolate the likelihood and query contract accordingly



Predictions for Instruction Count, Memory Accesses

Results for NAT, Bridge

Class	Description	
NAT1	Unconstrained traffic	
NAT2	Client packet, new flow	
NAT3	Existing flow	
NAT4	External, dropped packet	
BR1	Unconstrained traffic	
BR2	Broadcast traffic	
BR3	Unicast traffic	



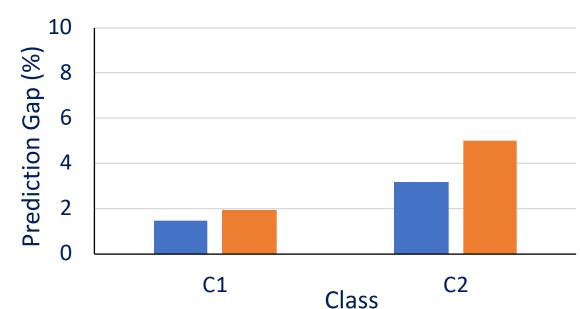
Bolt predicts IC & MA accurately, irrespective of NF/Traffic Class

Predictions for NF chains

$\,\circ\,$ NFs chained together

- Firewall drops packets with IP options
- Router Fast path (No IP options), Slow path (packets with IP options)

Class	Description	
C1	Packets with IP options	
C2	Packets without IP options	



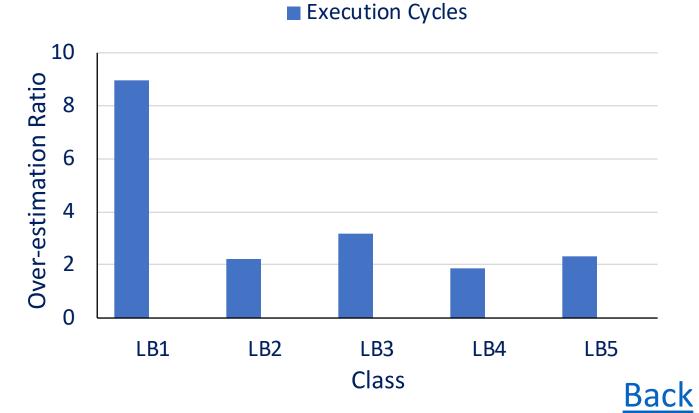
Instruction Count (IC) Memory Accesses (MA)



Predictions for Latency (Execution Cycles)

Results for Maglev-like Load Balancer

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LB5	Heartbeat packets	



46

9x for pathological traffic, 3x for typical traffic

Predictions for Execution Cycles

Results for LB,NAT, Bridge,LPM

NF+Class	Predicted Bound	Measured Cycles	Ratio
NAT1	591,948,908,371	65,217,699,390	9.08
NAT2	7,401	2,376	3.11
NAT3	5,142	1,789	2.87
NAT4	2,956	884	3.34
Br1	295,984,939,878	32, 383, 472, 634	9.14
Br2	7,329	2,013	3.64
Br3	7,383	1,808	4.08
LB1	591,969,879,756	66,062,284,173	8.96
LB2	5,299	2,386	2.22
LB3	8,108	2,541	3.19
LB4	4,300	2,310	1.86
LB5	4,837	2,079	2.33
LPM1	1,419	967	1.46
LPM2	1,015	545	1.86

Table 3: Accuracy of execution cycle performance contracts for multiple NFs and packet classes.



Full Blown Contract

Traffic Type	Instructions
Invalid packets (dropped)	$359 \cdot e + 80 \cdot e \cdot c + 38 \cdot e \cdot t + 425$
Known flows (forwarded)	$359 \cdot e + 30 \cdot c + 18 \cdot t + 80 \cdot e \cdot c + 38 \cdot e \cdot t + 1030$
New external flows (dropped)	$359 \cdot e + 30 \cdot c + 18 \cdot t + 80 \cdot e \cdot c + 38 \cdot e \cdot t + 528$
New internal flows; table full (dropped)	$359 \cdot e + 30 \cdot c + 18 \cdot t + 80 \cdot e \cdot c + 38 \cdot e \cdot t + 639$
New internal flows; table not full (forwarded)	$359 \cdot e + 30 \cdot c + 44 \cdot t + 80 \cdot e \cdot c + 38 \cdot e \cdot t + 1316$

Table 6: VigNAT performance contract. Instructions are described as a function of the number of expired flows (e) and the number of hash collisions (c) and bucket traversals (t) incurred in the hash table.

